

¿POR QUÉ TIPOS DE DATO ALGEBRAICOS?



Sum Type y búsqueda de patrones sobre ellos



Run Reset

EDIT ON
ScalaFiddle

RESULT

```
sealed trait Triple
case class L(l: Boolean) extends Triple
case class M(m: Boolean) extends Triple
case class R(r: Int) extends Triple

val r : Triple = L(true)

val result = r match {
  case L(true) => "L_True"
  case M(false) => "M_False"
  case R(1) => "R_1"
  case _ => "Perl6"
```

run ▶

open in repl.it

main.hs

```
1 data Triple = L Bool | M Bool | R Int
2 f :: Triple -> String
3 f x = case x of L True -> "L True"
```

GHCi, version 8.6.3

Product Type y búsqueda de patrones



Run Reset

EDIT ON
ScalaFiddle

```
case class Pair[A, B] (a: A, b: B)

def processPair[A,B](pair: Pair[A,B]) = {
  pair match {
    case Pair(fst: Int, snd :Int) => "Pair (Int, Int) "
    case Pair(fst: Int, snd :String) => "Pair (Int, String)"
    case Pair(fst: Boolean, snd :Boolean ) => "Pair (Bool, Bool)"
    case _ => "Pair of something else"
  }
}

// println(processPair(Pair(1,"2")))
```

RESULT

run ▶

open in repl.it

main.hs

```
1 data Pair a b = P a b deriving Show
2 pairBoolBool :: Pair Bool Bool
3 pairBoolBool = P True True
```

GHCi, version 8.6.3

Exponential type



Run Reset

EDIT ON
ScalaFiddle

```
type BoolToBoolET = Boolean => Boolean
val notF : BoolToBoolET = (b: Boolean) => ! b

// Additional :: Partial Type Application via alias
type T[A] = Option[Map[Int, A]]
val t: T[String] = Some(Map(1 -> "abc", 2 -> "xyz"))
```

RESULT

run ▶

open in repl.it

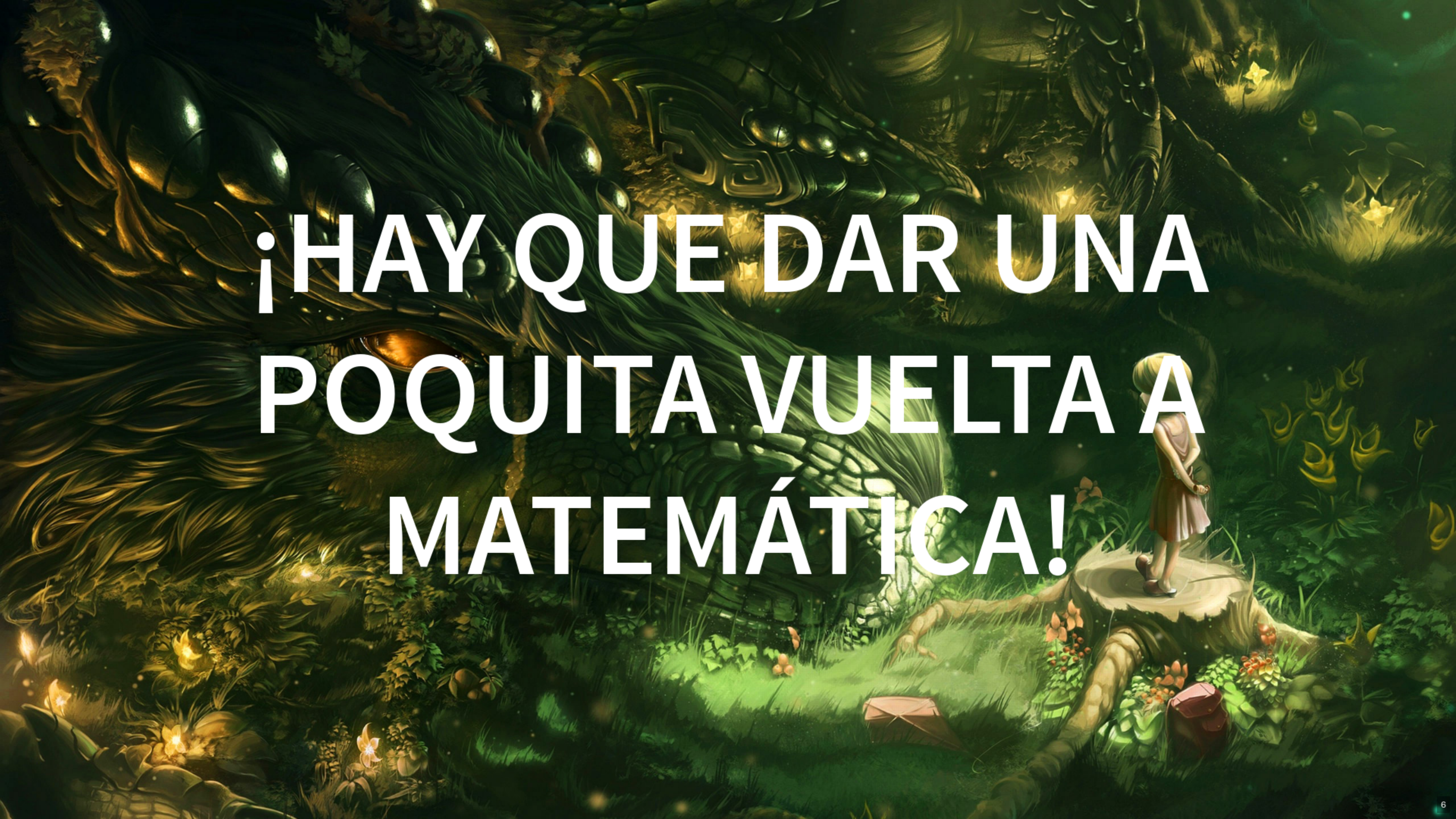
main.hs

```
1 type BoolToBool = Bool -> String
2 notAsString :: BoolToBool
3 notAsString True = "Fasle"
```

GHCi, version 8.6.3

PERO...

¿POR QUÉ SON
ALGEBRAICOS?

A vibrant, fantastical illustration of a magical forest. In the foreground, a young girl with blonde hair, wearing a pink dress, stands on a large, mossy tree stump. She is looking down at a small, glowing red object on the ground. The forest is lush with green foliage, and the ground is covered in moss and small, glowing flowers. In the background, a large, gnarled tree trunk is visible, and the scene is illuminated by warm, golden light. The overall atmosphere is whimsical and enchanting.

**¡HAY QUE DAR UNA
POQUITA VUELTA A
MATEMÁTICA!**

CADA **ADT** LLEVA SU OPERADOR PARA REALIZAR SU ÁLGEBRA

$$a \mid b = |a| + |b|$$

$$(a, b) = |a| * |b|$$

$$a \rightarrow b = |b| ^ |a|$$

PARTE UNO

LOS TIPOS Y SUS CARDINALIDADES

```
|Void|    = 0
|()|      = 1
|Bool|    = 2
-- De aquí ya podemos realizar lo que siguiente
type Product = |(Bool, Bool)|      = ???
type Sum = |Either Bool Bool|      = ???
type Exponential = |(Bool -> Bool)| = ???
```

La correspondencia unívoca tipos A y B con las mismas cardinalidades son isomorfismos
TODO absurd :: Ex falso sequitur quodlibet

ISOMORPHISMO CURRY-HOWARD

Álgebra	ADT
$a + b \Leftrightarrow \Sigma$	Either a b
$a * b$	(a, b)
$b^a \Leftrightarrow \prod b$	$a \rightarrow b$
1	()
0	void

PARTE DOS - DADA LA DEMOSTRACIÓN

ISOMORFISMO ENTRE ENTRE MATEMÁTICA Y TIPOS

```
(a * b)^c = a^c * b^c
a * (b + c) = a * b + a * c
-- not that easy
(a^b)^c = a ^ (b * c)
-- TODO show maxBound
-- TODO show Data.Void
-- curry/uncurry
```


LA FORMA CANÓNICA DE ADT

$$t = \sum \prod t_m, n_m n$$

```
-- canónicas ::  
Either a b  
a -> b  
Either a (Either b, (c,d))  
  
-- no canónicas ::  
(a, Either b c)  
(a, Bool)
```

LAS VARIANTES DEL TIPOS

*“A Profunctor is a Contravariant
Functor on its first type parameter and
a Functor on its second type parameter.”*

data Variance = Covariant | Contravariant | Invariant
import Data.Functor

*A Variance of type $T\ a$ is fully specified
by positive, negative (or mix of both)
type variable position.*

```
newtype F1 a = F1 ( Int -> a ) -- + => + Covariant
newtype F2 a = F2 ( a -> a ) -- +/- => +/- Invariant
newtype F3 a = F3 ( (Int -> a) -> Int ) -- Contravariant - * + =>
newtype F4 a = F4 ( (a -> Int) -> Int ) -- Covariant - * - => +
```

SCALA -> (HASKELL, IDRIS, PERL 6)

=> PATTERN MATCHING - POWER OF THE TYPES

=> TYPE CLASSES VS. INTERFACES

=> KINDS - AS A LABELS FOR THE TYPES

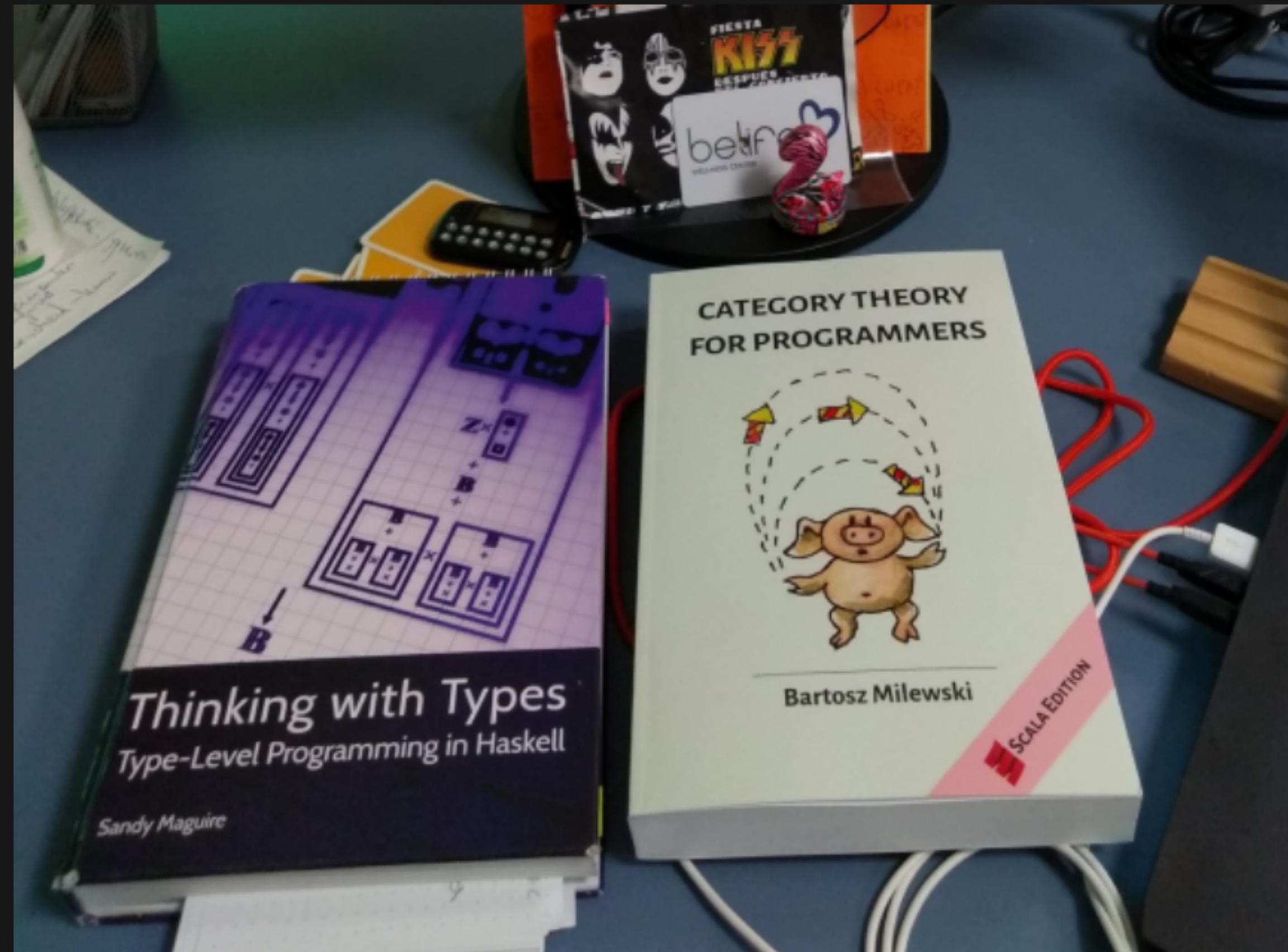
=> GADTS - GENERALIZED ADT

**=> PARTIAL TYPE APPLICATION/TYPE LEVEL
PROGRAMMING**

=> A BETTER SCALA REPL [AMMONITE.IO](https://ammonite.io)

=> TYPELEVEL [CATS LIBRARY](https://typelevel.org/cats/)

=> SCALAZ [SCALAZ-ZIO](https://scalaz.github.io/scalaz/)



¿Why?

Because writing like this we programming by giving
proof.

¡Think about it!

!!EXTRA!! GADTS



```
role Maybe[::A] { }  
sub just(::A $x --> Maybe[::A]) { class Just does Maybe[::A] { has $.V = $x; }.new }  
sub nothing(Any:U $t --> Maybe[$t]) { class Nothing does Maybe[$t] { }.new }
```



```
data MyMaybe a where  
  MyJust      :: a -> MyMaybe a  
  MyNothing :: MyMaybe a
```

¡MUCHAS GRACIAS!



THANKS!